

PLANAR MAPS

(Randomness & Enumeration)

↑ just a bit

↑ mostly

guillaume.chapuy@irif.fr

Introduction to very active topic.

Will not talk about:

- higher genus
- Mathematical physics
- Matrix integrals
- Rep. theory / alg. combinatorics
- ...

Mostly: counting maps by

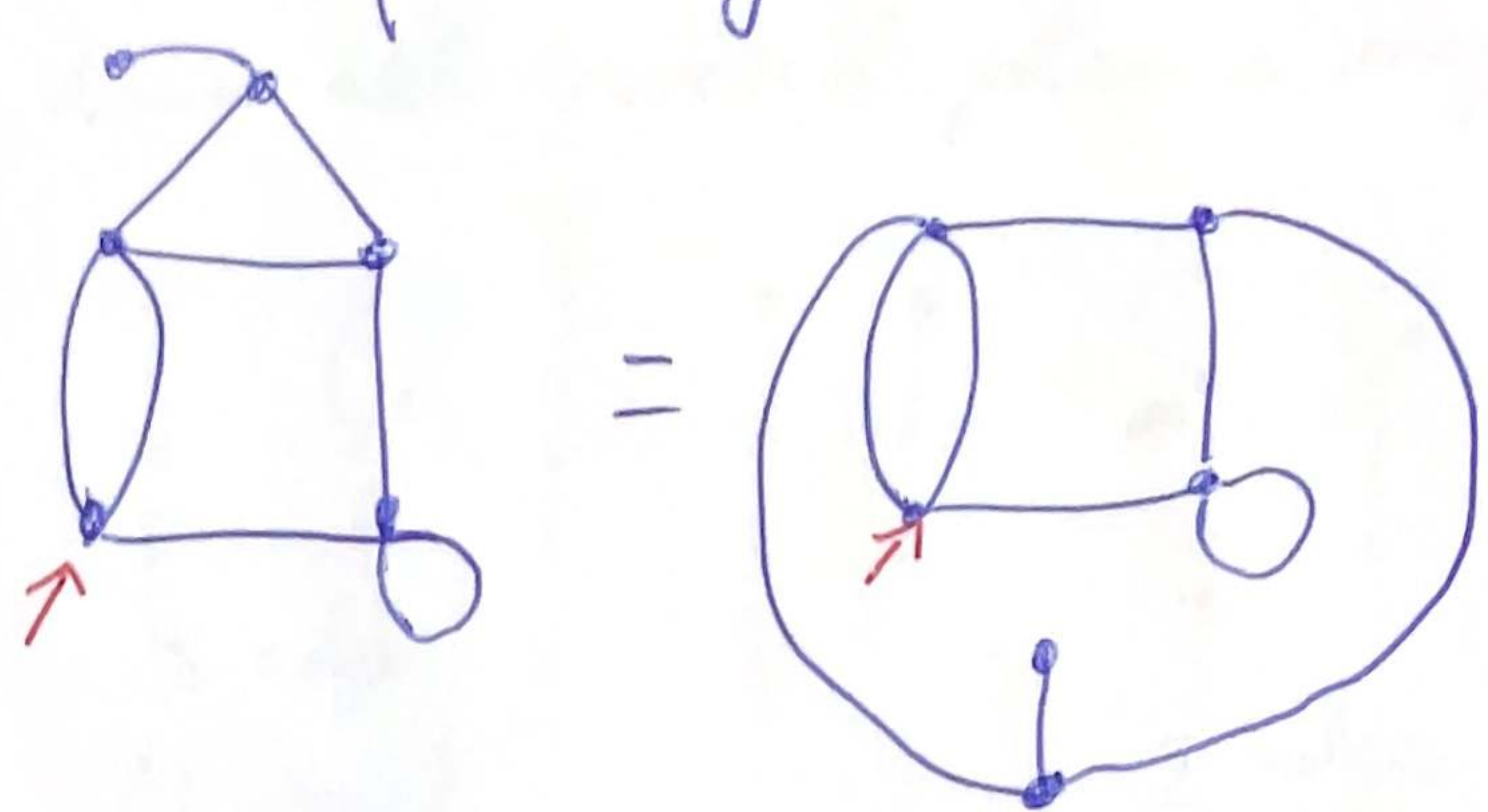
- 1/ Functional equations
- 2/ Surgery (= bijections)

- basic consequences in probability

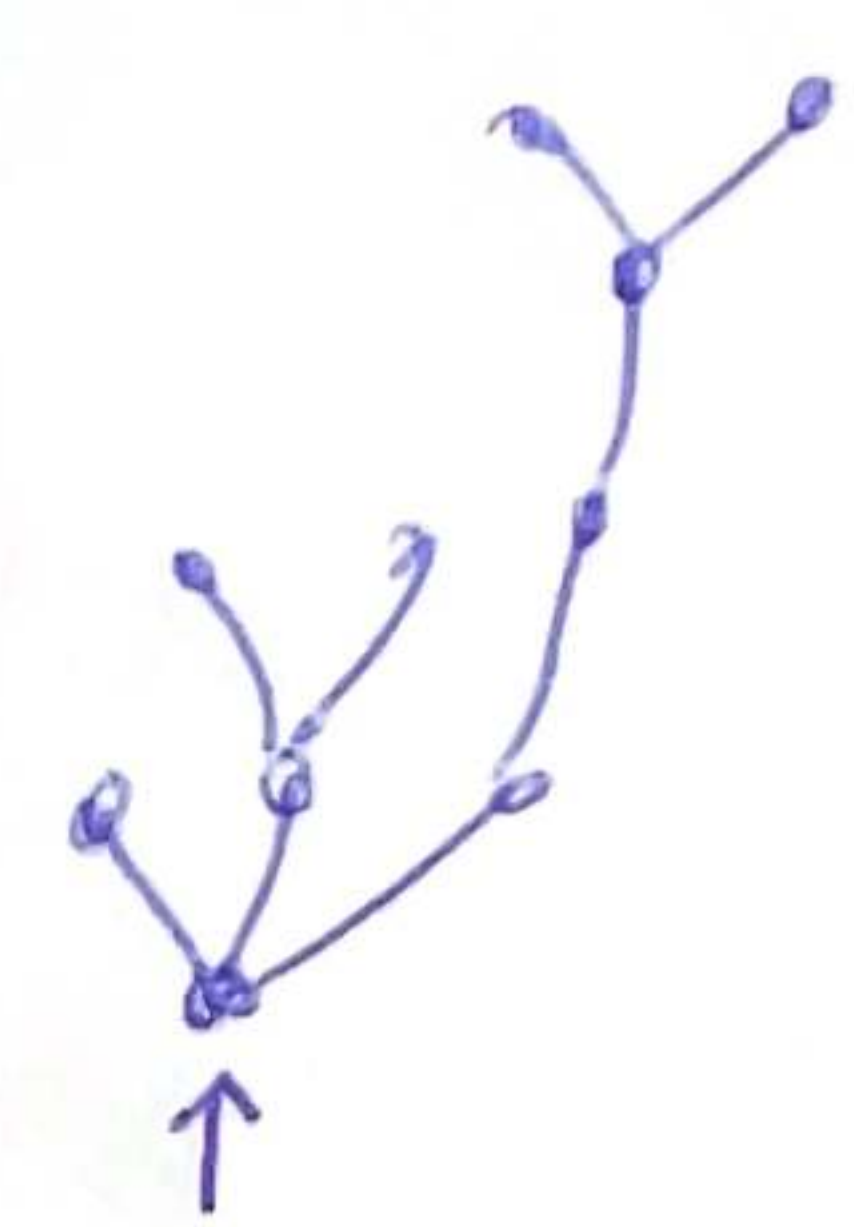
Refs: see course page on my website

SFB Combinatorics
& Probability
Summer School
Potsdam am Wölkensee
Austria, Sept 2025

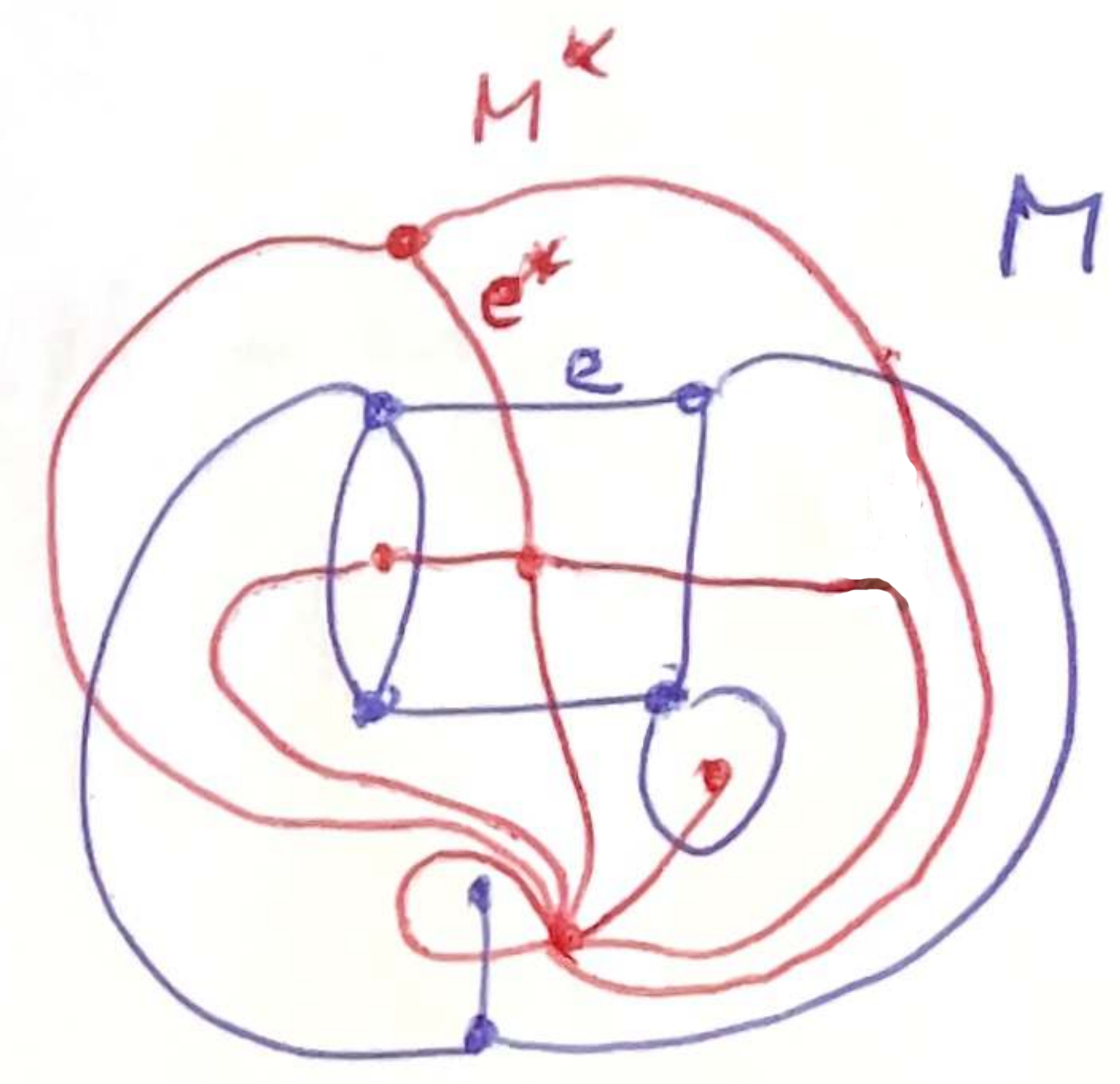
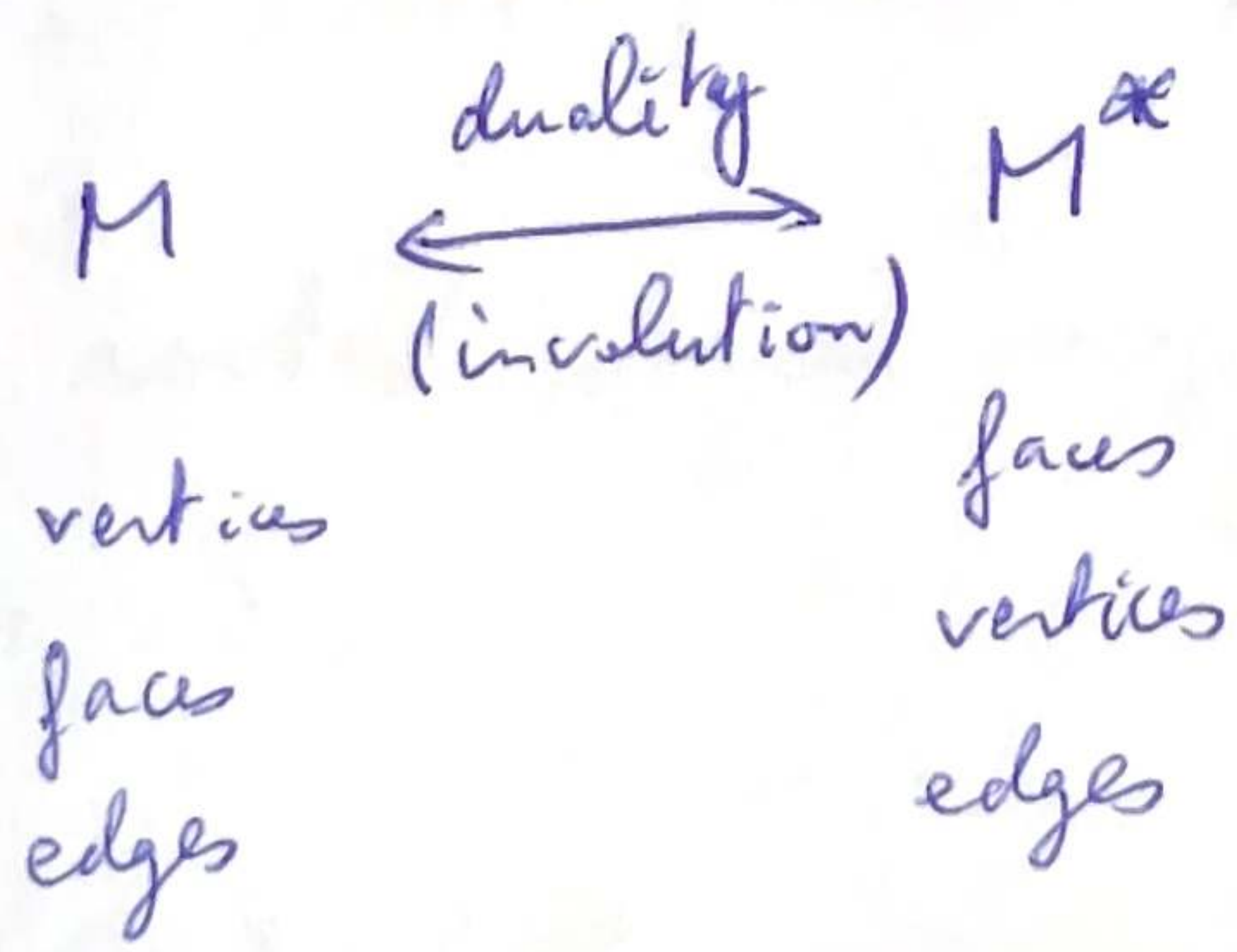
Rooted map = a distinguished corner
(usually drawn on the "exterior" face)



Note: rooted (planar) map with one face
= plane rooted tree (Catalan numbers!)



Dual map:



Dual submap: M map, $S \subseteq E(M)$
 $S^* = \{e^*, e \in S\} \subseteq E(M^*)$

Fact: S^* induces a connected graph on $V(M^*)$ iff S has no cycle
(Jordan lemma)

Cor1: S spanning tree of $M \iff S^*$ spanning tree of M^*

Cor2: Euler's formula $V + F = E + 2$

↳ pf: let S be a spanning tree of M .

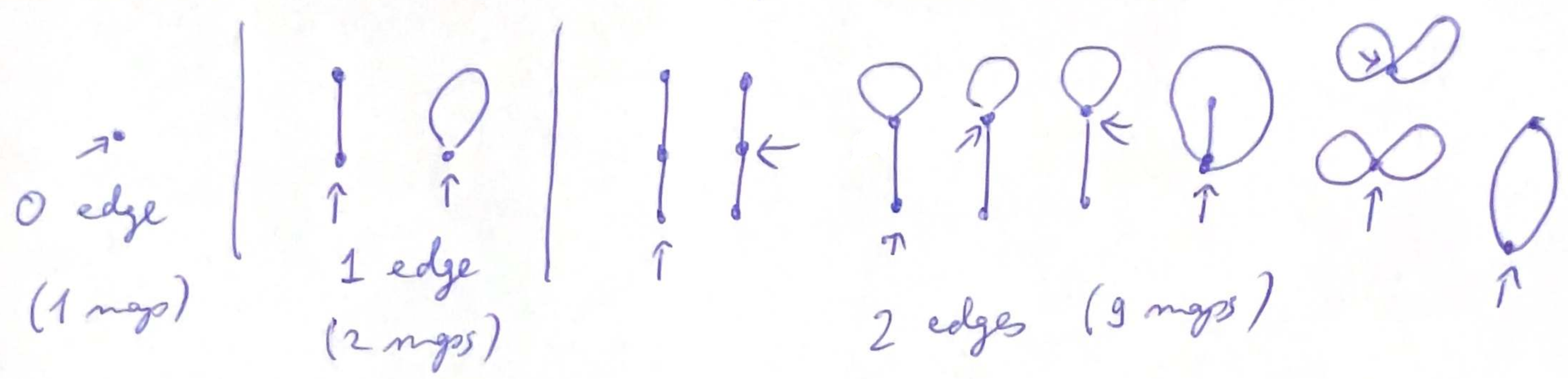
$$1 + E(S) = V(M) = V$$

$$1 + E(S^*) = V(M^*) = F$$

$$\Rightarrow 2 + E = V + F \quad \square$$

II / Counting

Let's draw all rooted planar maps up to 2 edges:



Warm-up (see exercise session)

The number of tree rooted maps (maps equipped with a spanning tree) is $Cat(n) Cat(n+1)$.

Let $m_n =$ nb of rooted planar maps with n edges

$$(m_n)_{n \geq 0} = 1, 2, 9, 54, \dots$$

Thm (Tutte, 1960s) (the real thing)

$$m_n = \frac{2}{n+2} 3^n Cat(n)$$

Equivalently $M(z) = \sum m_n z^n$ is given by
$$\begin{cases} M = R - 3R^3 \\ R = 1 + 3zR^2 \end{cases}$$

COR: $m_n \sim \frac{2}{\sqrt{\pi}} n^{5/2} 12^n$ $\rightarrow$ "planar map exponent"

Proof by Tutte equation...

Tutte equation

Let us remove the root edge in a map.

Two cases $\rightarrow$ disconnecting or not disconnecting.



Let $M(z, u) := \sum_{\text{rooted planar maps}} z^{\text{edges}} u^{\text{degree of root face}}$

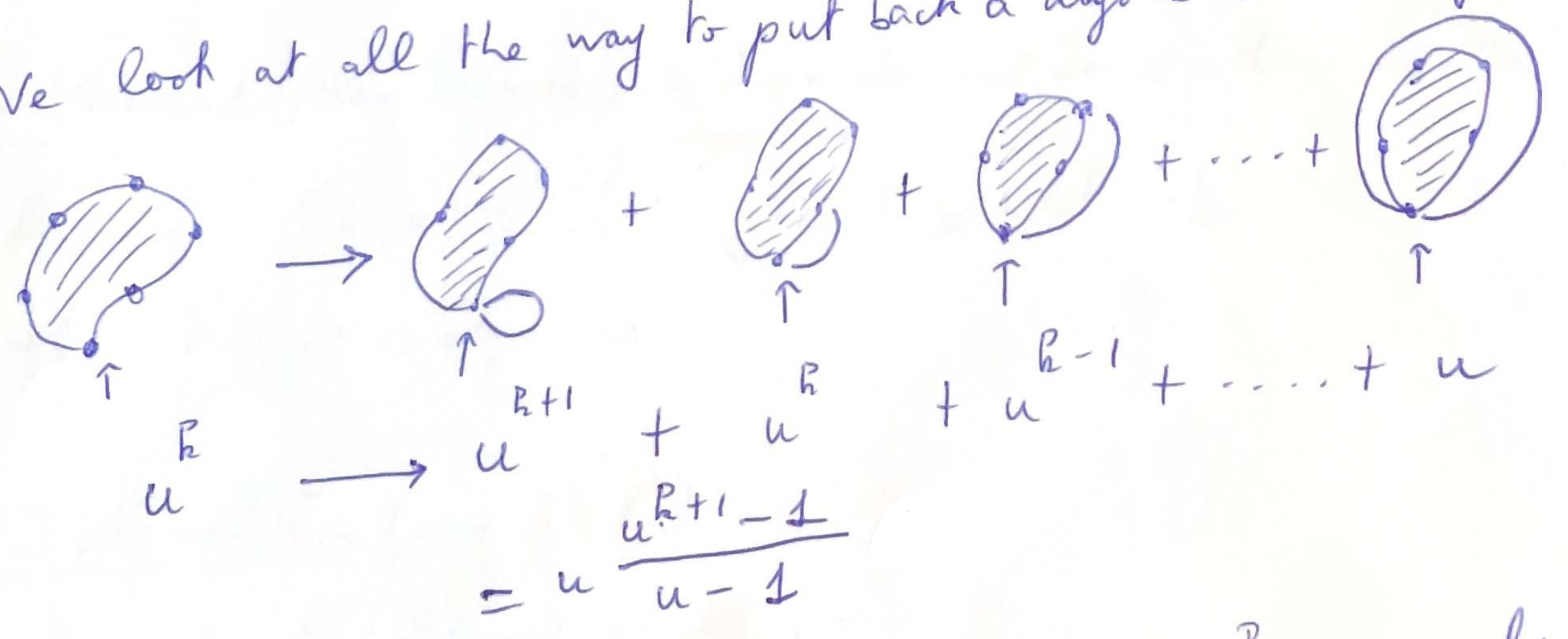
"catalytic variable"

Fact: we can write an equation for $M(z, u)$!

Contribution of (i): $z u^2 M(z, u)^2$

(ii): ?

We look at all the way to put back a diagonal in the root face:



Let $\Delta: A(u) \mapsto \frac{A(u) - A(1)}{u - 1}$. Then we have proved:

(*) $N = 1 + z u^2 N^2 + z u \Delta(u N)$

Tutte equation!

$$N(z, u) = 1 + z u^2 N(z, u)^2 + z u \frac{u N(z, u) - N(z, 1)}{u - 1}$$

Rem. - (*) determines a unique $N(z, u) \in \mathbb{C}[u][[z]]$

- even more: see this equation as $P(N(z, u), N(z, 1), z, u) = 0$

with a polynomial P , then $P(A(z, u), a(z), z, u)$
has a unique solution in $\mathbb{C}[u][[z]] \times \mathbb{C}[[z]] \ni (A, a)$
and it is such that $A(z, 1) = a(z)$

Method 1 - guess $N(z, 1) = \sum \frac{2 \cdot 3^n}{n+1} \text{Cat}(n) z^n$

(Tutte) - solve for $N(z, u)$

- check that $N(z, u) \in \mathbb{C}[u][[z]]$

Thm (Popescu, 1980s)

If an equation of the form $P(A_1(z, u), \dots, A_k(z, u), a_1(z), \dots, a_\ell(z)) = 0$
has a unique solution in $\mathbb{C}[u][[z]]^k \times \mathbb{C}[[z]]^\ell$,
then each of these series is algebraic

Thm (Bousquet-Mélou Jehanne, slightly weaker but comes with an algorithm!)

Assume $F(z, u) = P_0(u) + z \overset{\text{polynomial}}{P_1}(F(z, u); \Delta F(z, u); \dots, (\Delta^R F)(z, u))$
then $F(z, u)$ is algebraic

BMJ algorithm (case $k=1$)

Assume $E(A(z, u), a(z), z, u) = 0$

Look for $u=U(z)$ such that $(\partial_1 E)(A(z, u), a(z), z, u) = 0$

These two equations imply $(\partial_4 E)(A(z, u), a(z), z, u) = 0$

(since $E=0 \Rightarrow A'(u)\partial_1 E + \partial_4 E = 0$)

→ we have a system
$$\begin{cases} E(A, a, z, u) = 0 \\ (\partial_z E)(A, a, z, u) = 0 \\ (\partial_u E)(A, a, z, u) = 0 \end{cases}$$

with three unknowns $(A = A(z, u(z)), a = a(z), u = u(z))$.

↳ demo with Maple! (but doable by hand of course)

→ Tutte's formula! $\square$

Complement: remarkable very recent result:

Thm [Contat-Cuic 2025+, $R=1$: C/Dmitry-Ney-Yu...]

A BMS-type equation with nonnegative coefficient,
under mild hypotheses (including non-linearity!)
always has an $-\frac{5}{2}$ counting exponent:

$$[z^n] F(z, 1) \underset{n \rightarrow \infty}{\sim} C n^{-5/2} K^n$$

(great result! Follows from ^{deep} general classification
of some self-similar random trees).

III Bijections

rooted maps $\mathcal{M}_n$ with n edges $\rightarrow |\mathcal{M}_n| = \frac{2}{n+2} 3^n \text{Cat}(n) \quad ???!$

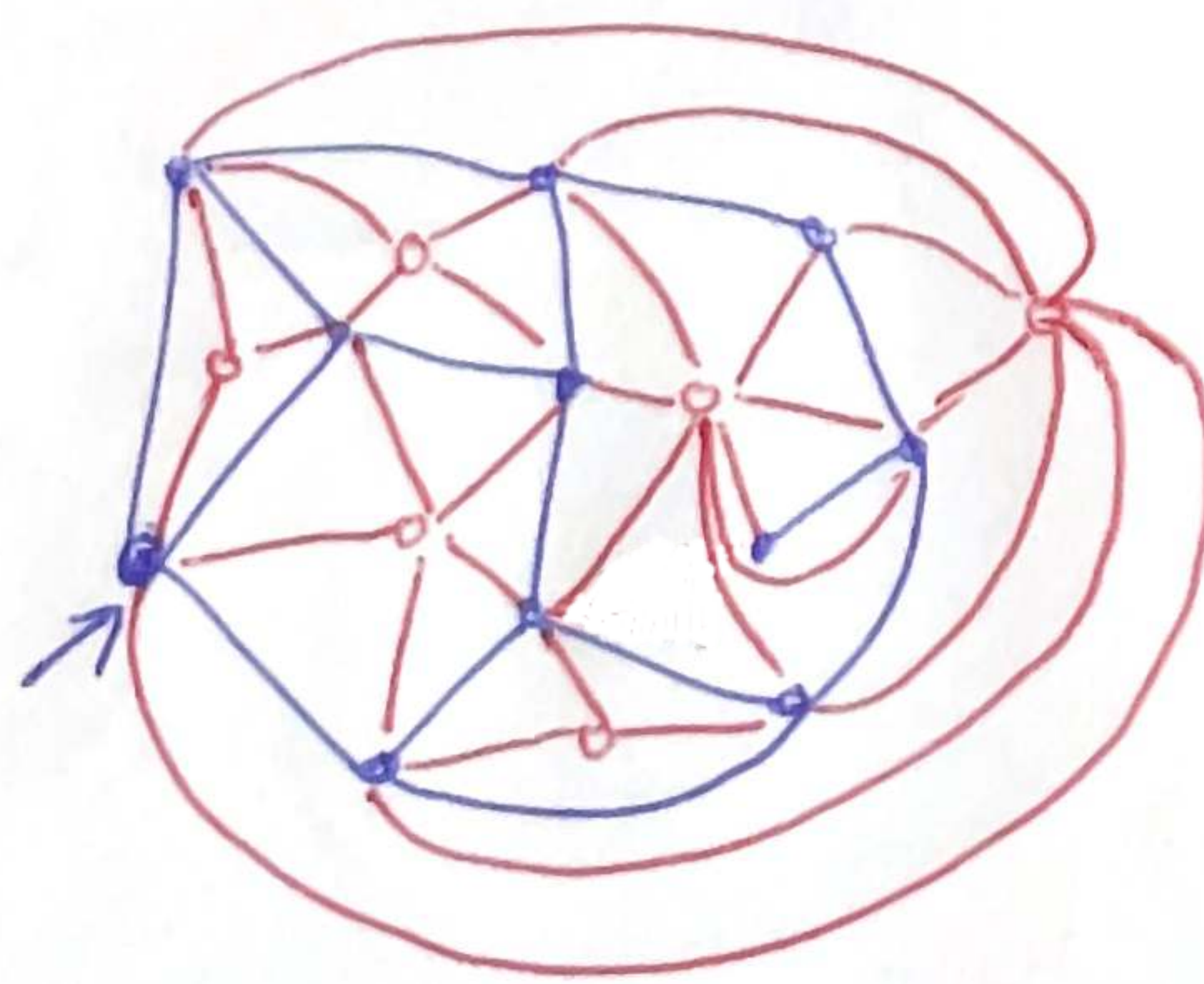
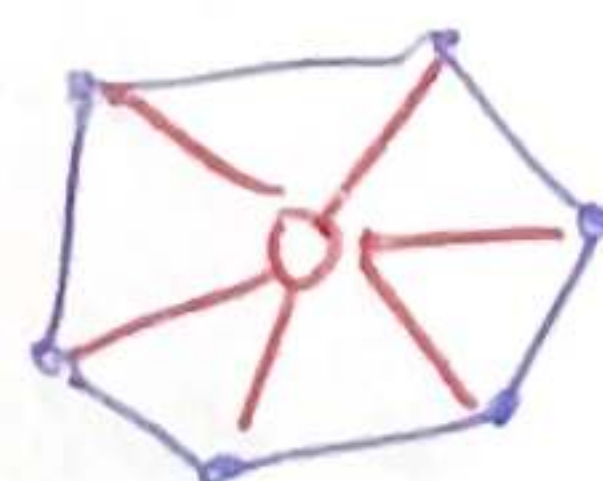
quadrangulation = map with all faces of degree 4.

$$\mathcal{Q}_n = \{\text{rooted quadrangulations with } n \text{ faces}\}.$$

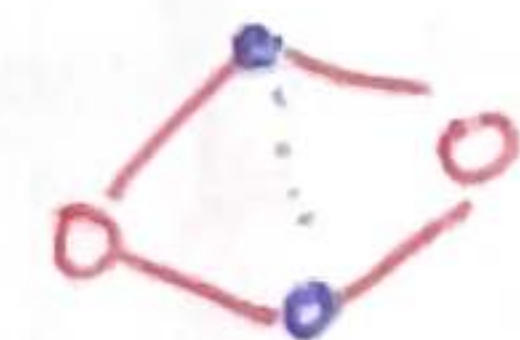
Fact: if $q \in \mathcal{Q}_n$, q has $n+2$ vertices
(pf: $v+f=e+2$, $4f=2e$)

Fact: if $q \in \mathcal{Q}_n$, then q is bipartite
(pf: see exercises!)

Tutte bijection: $\mathcal{M}_n \xleftrightarrow{\text{bij}} \mathcal{Q}_n$



(forget blue edges!)



Conversely:



let $\mathcal{Q}_n^\bullet = \{(q, v), q \in \mathcal{Q}_n, v \in V(q)\}$

$$|\mathcal{Q}_n^\bullet| = (n+2) |\mathcal{Q}_n| = 2 \cdot 3^n \cdot \text{Cat}(n) \quad ???!$$

Let us prove this combinatorially!

Cori - Vauquelin - Schaeffer - Bouttier - Guillemet
trees $\cong$ slices

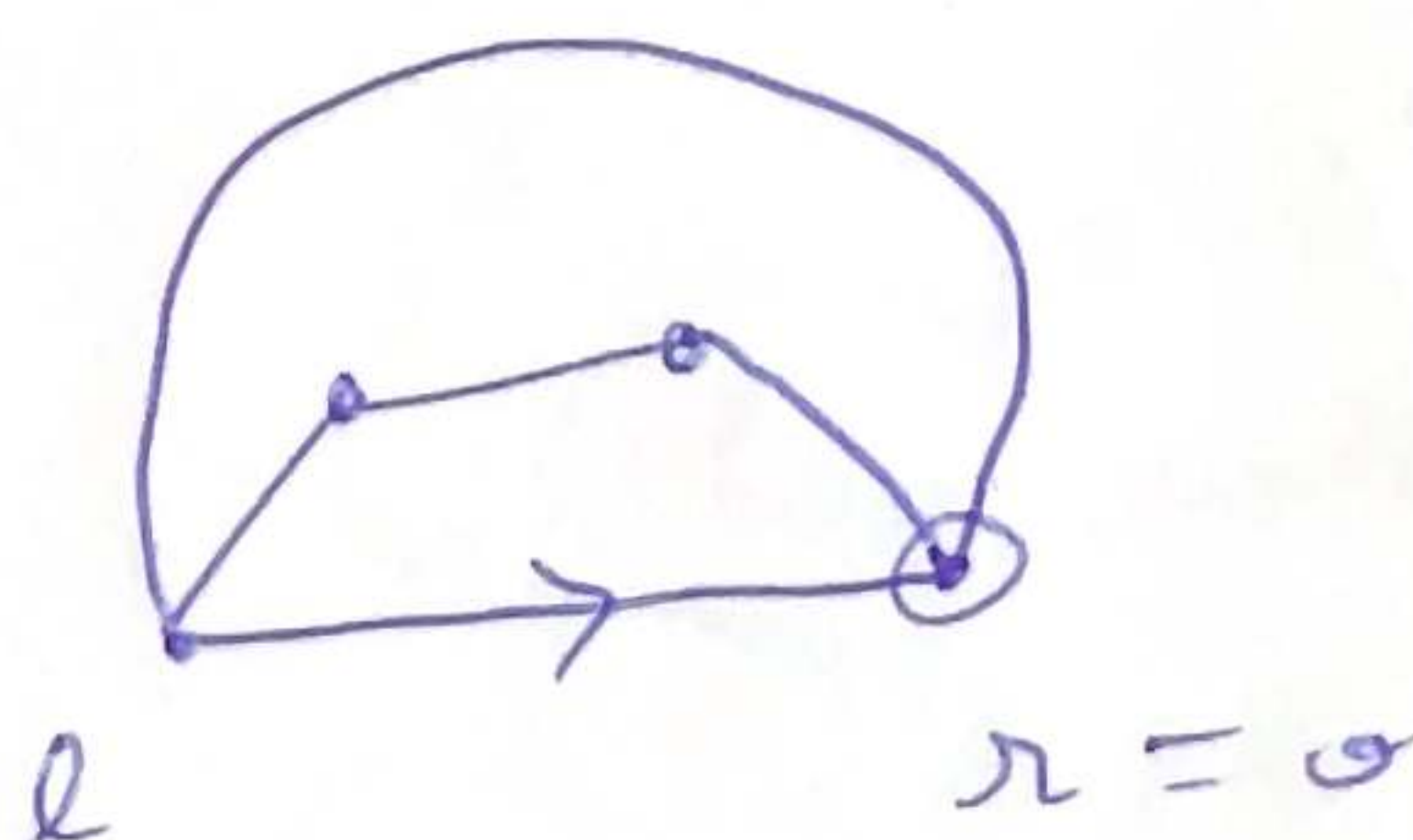
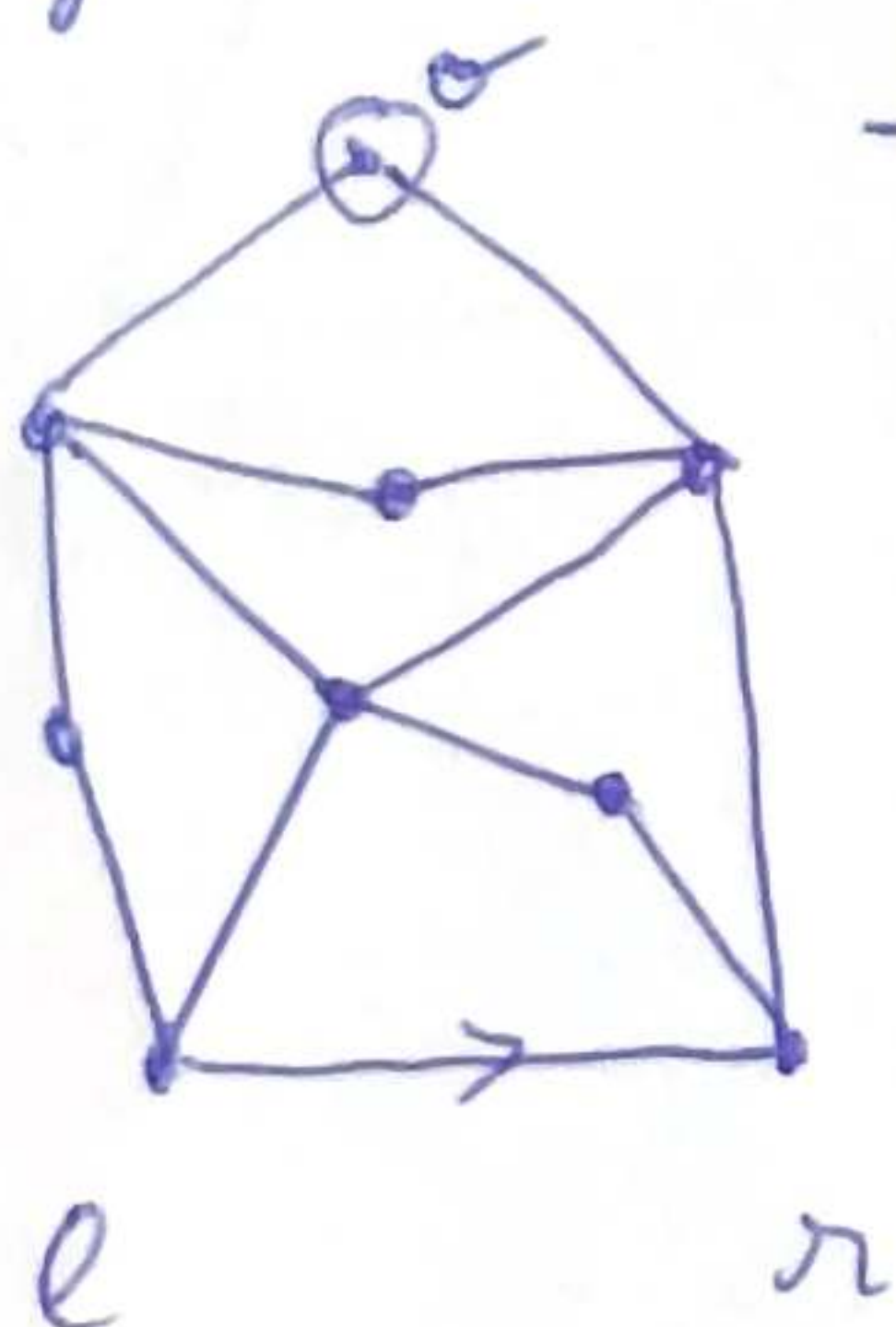
An elementary slice is a planar map with:

- all faces quadrangular except external one
- a "base" oriented edge $l \rightarrow r$ in the external face
- an "apex" ~~corner~~ σ in the external face

such that:

- left boundary is a geodesic $l \rightsquigarrow \sigma$
- right boundary is the unique geodesic $r \rightsquigarrow \sigma$
- apex is the only vertex belonging to both boundaries

ex :

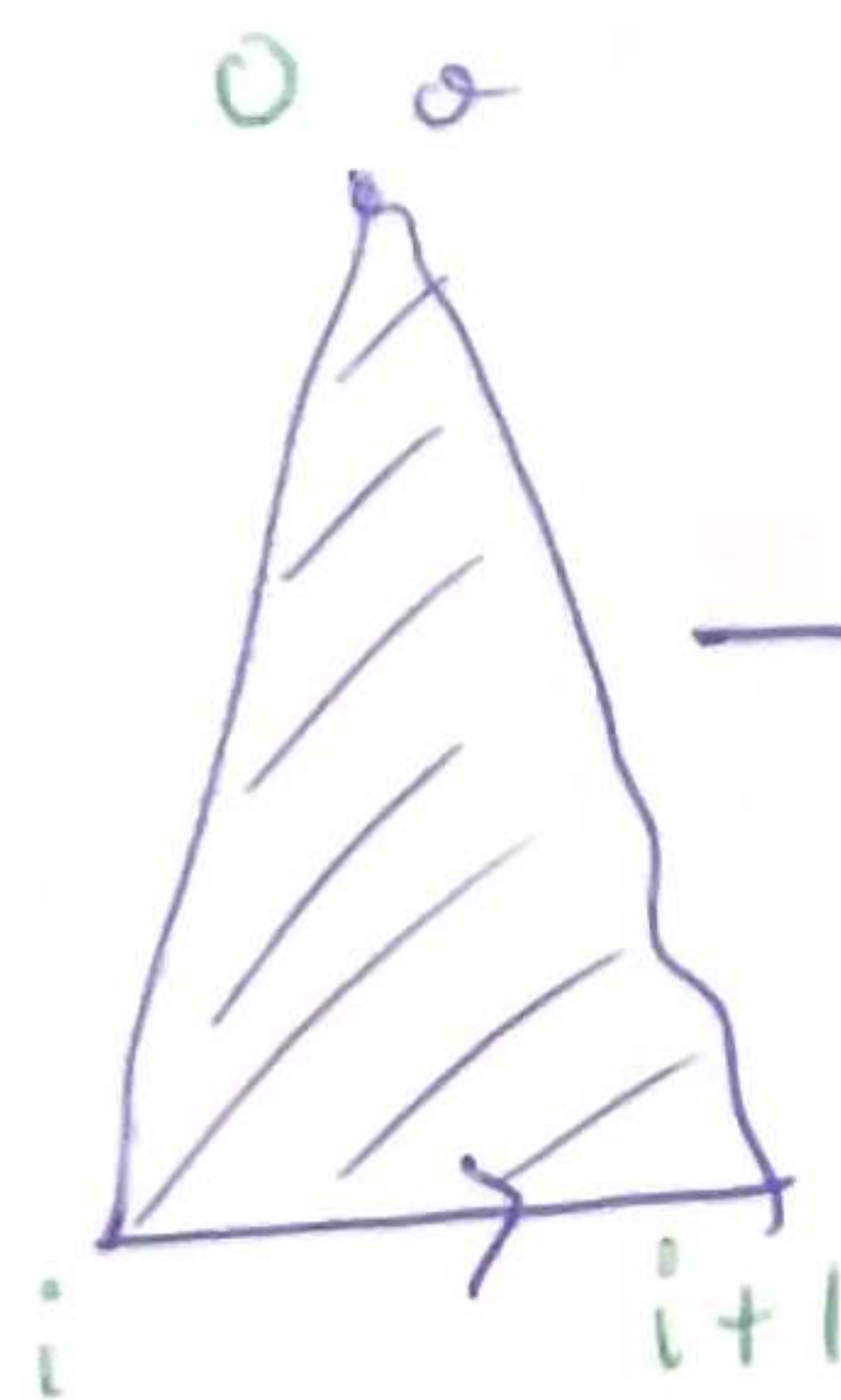


We label vertices by their distance to σ .

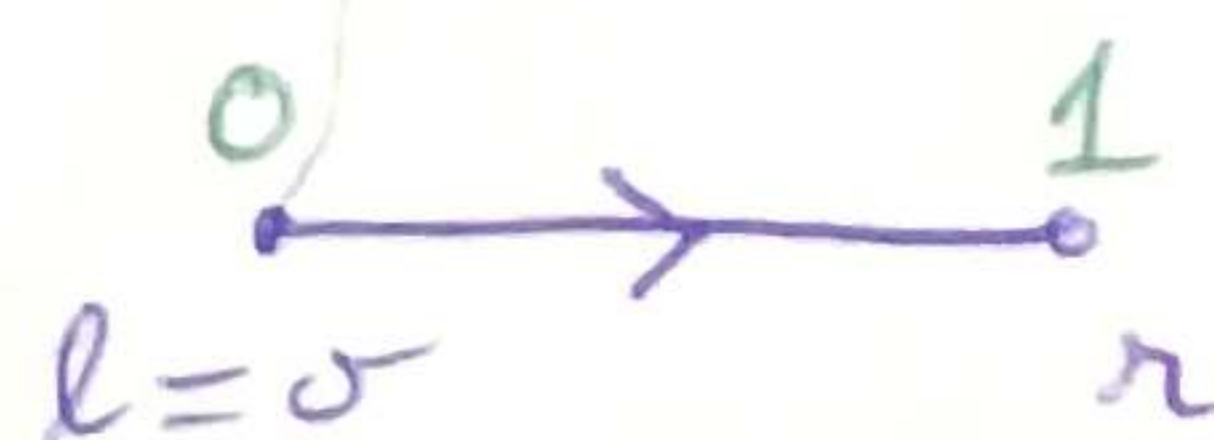
Only two cases:



or

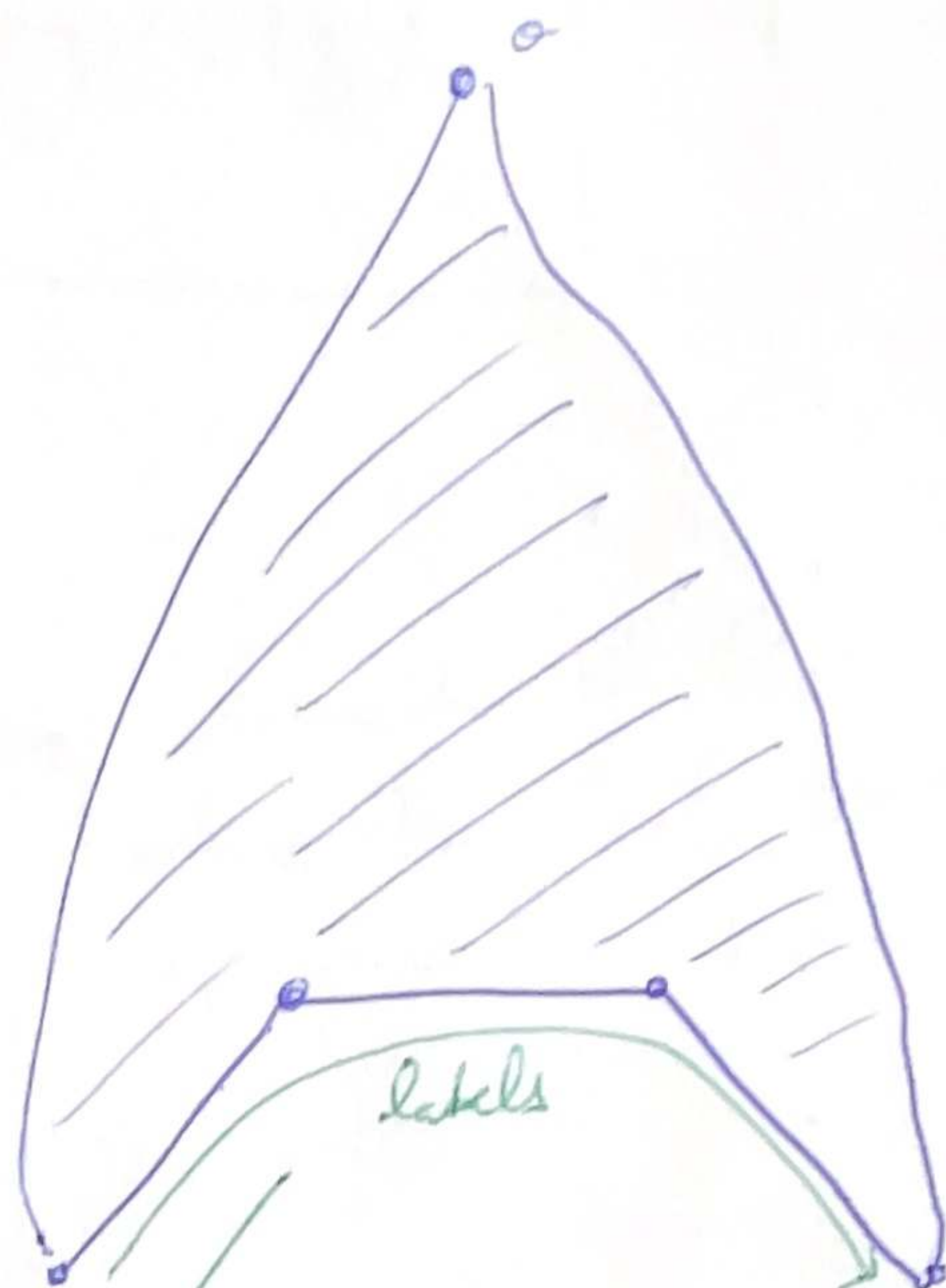
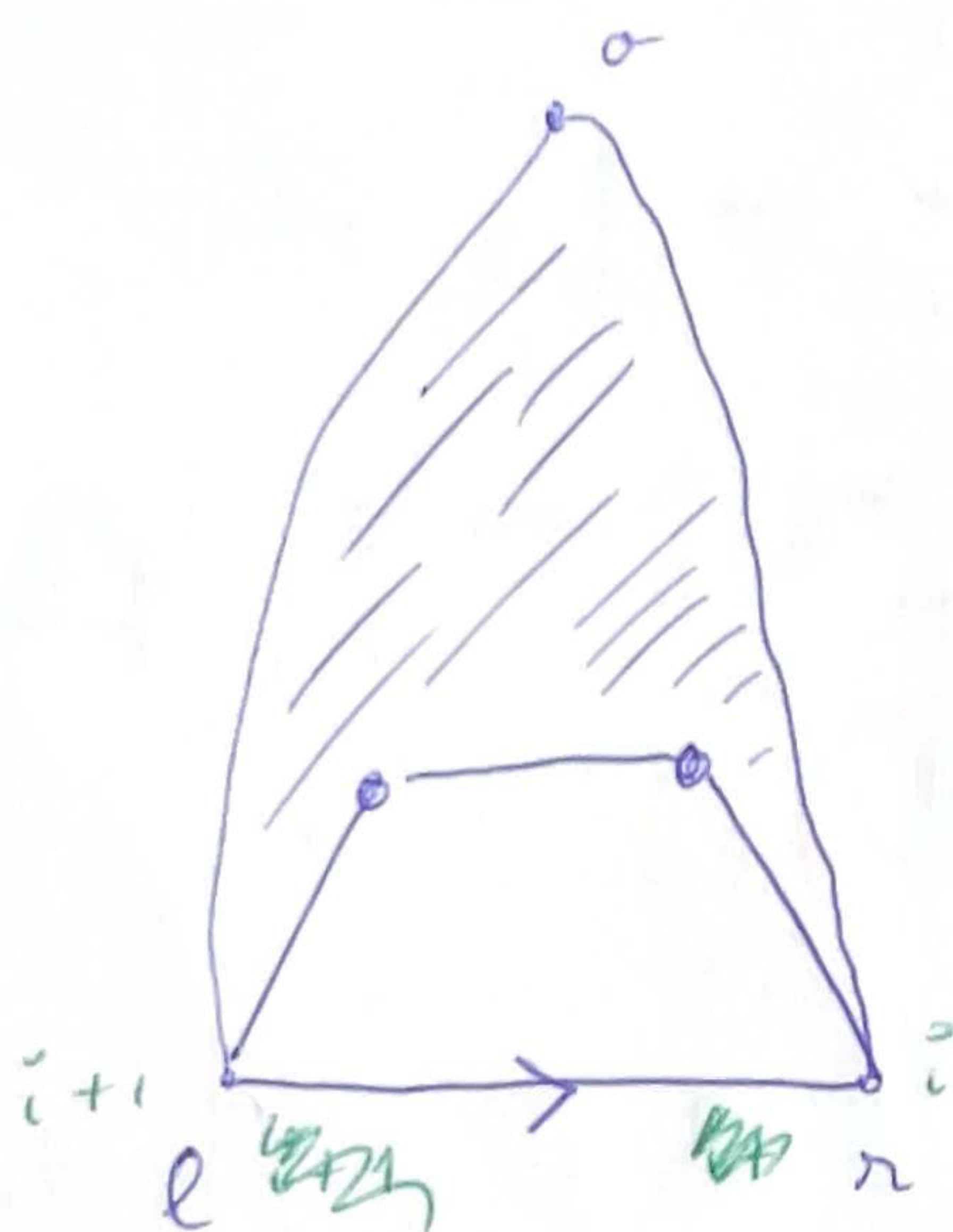


only in the one degenerate case



Let $R(z) =$ generating function for the first case.

Let us explore a slice!

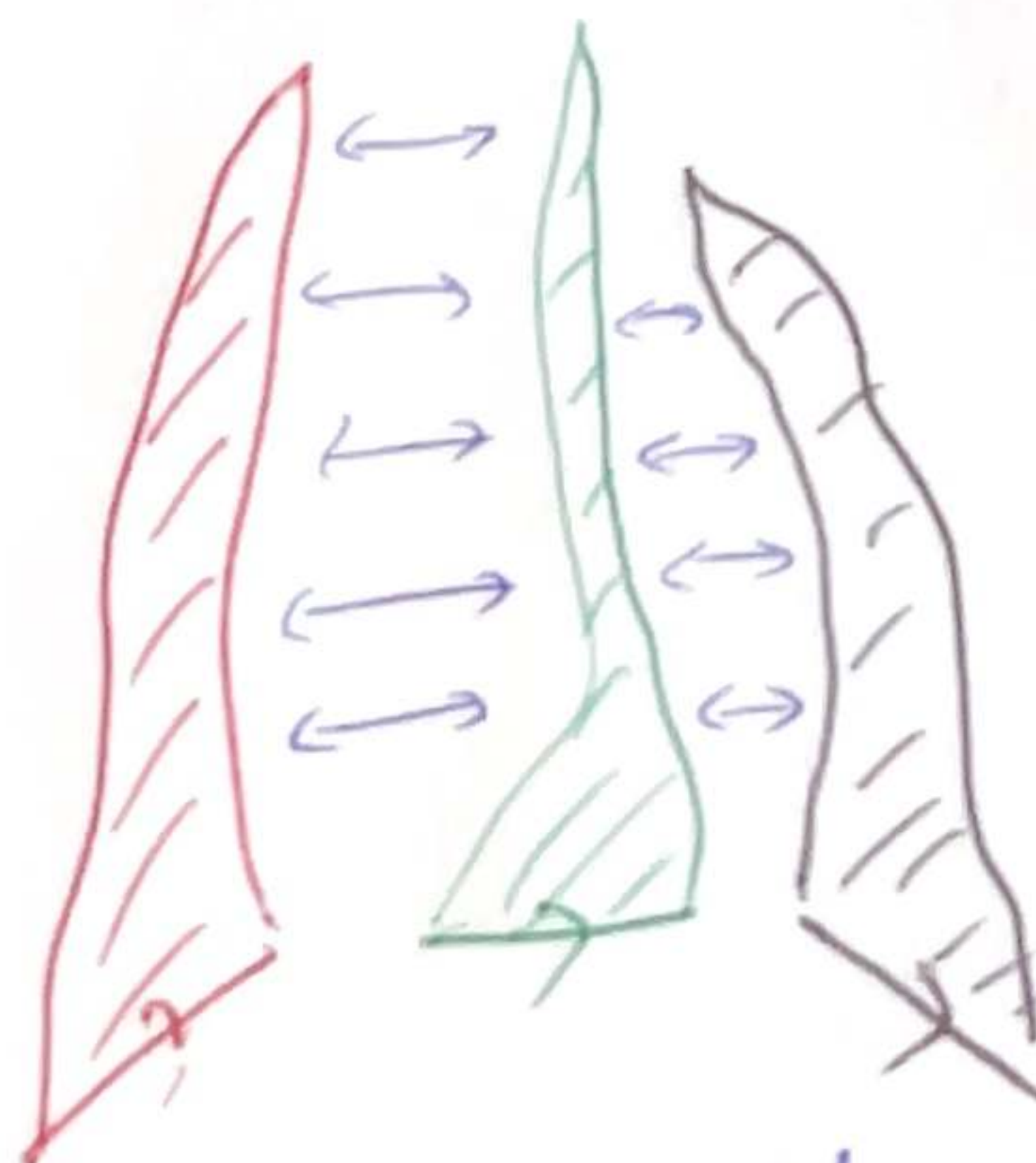
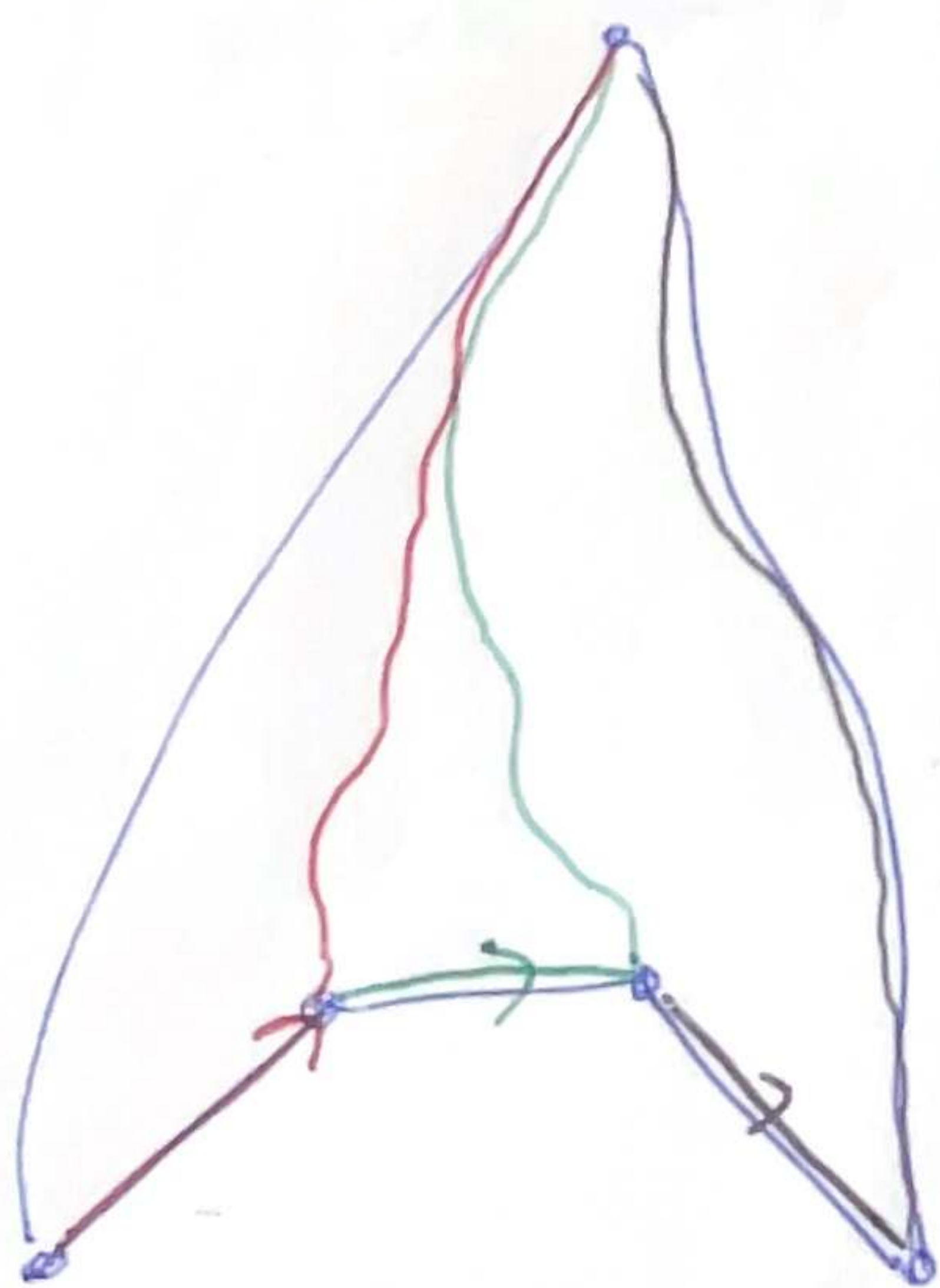


Three possible label sequences.

(*) $\left\{ \begin{array}{l} i+1 \# i+2 \# i+1 \# i \\ \text{or } i+1 \# i \# i+1 \# i \\ \text{or } i+1 \# i \# i-1 \# i \end{array} \right.$

$\hookrightarrow$ in each case,
 - one ascent
 - two descents

Further explore
 the left-most geodesic
 starting from each of the
 three edges.



Three slices!!

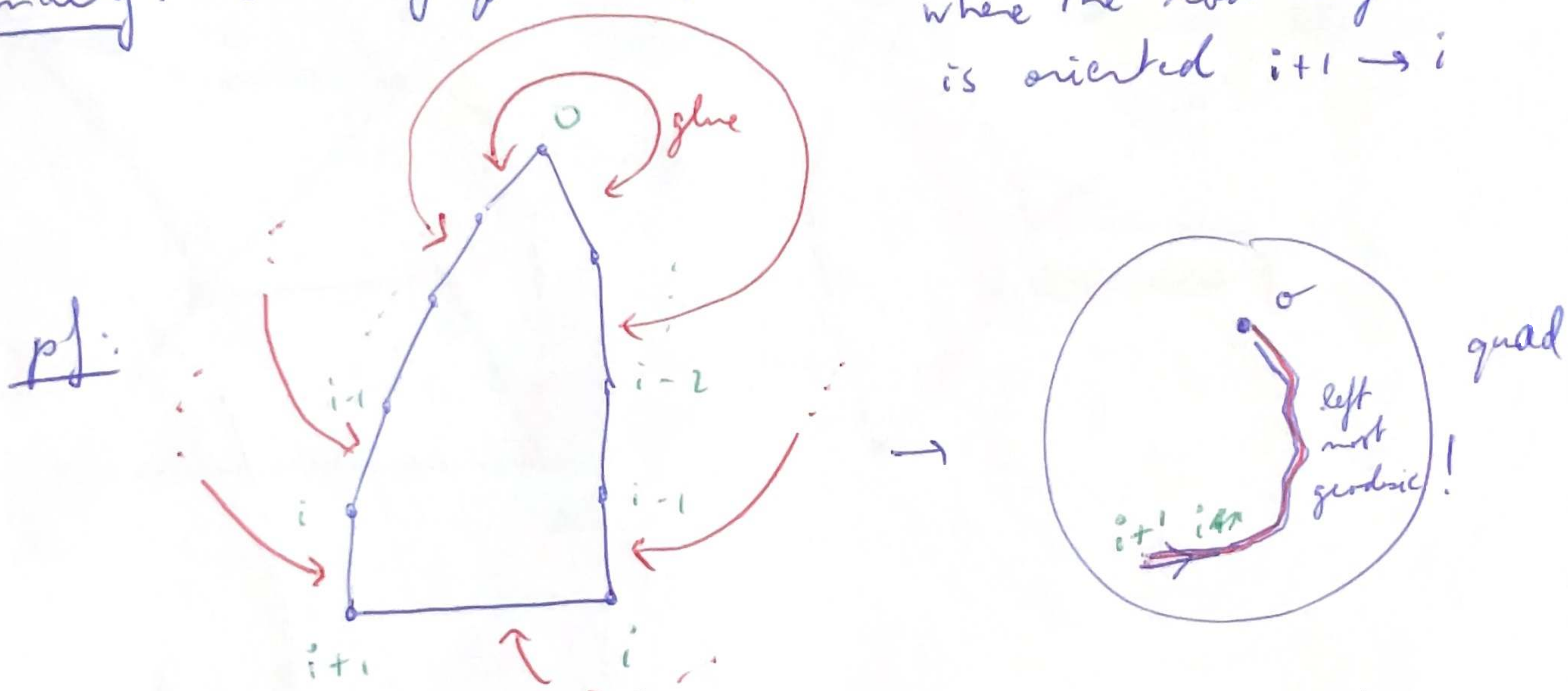
$\hookrightarrow \left. \begin{array}{l} \text{one trivial } (i \rightarrow i+1) \\ \text{two of type "R(3)"} \end{array} \right\} \text{ by (*)}$

Conclusion:

$$R(z) = 1 + \underbrace{(3)}_3 R(z)^2 !$$

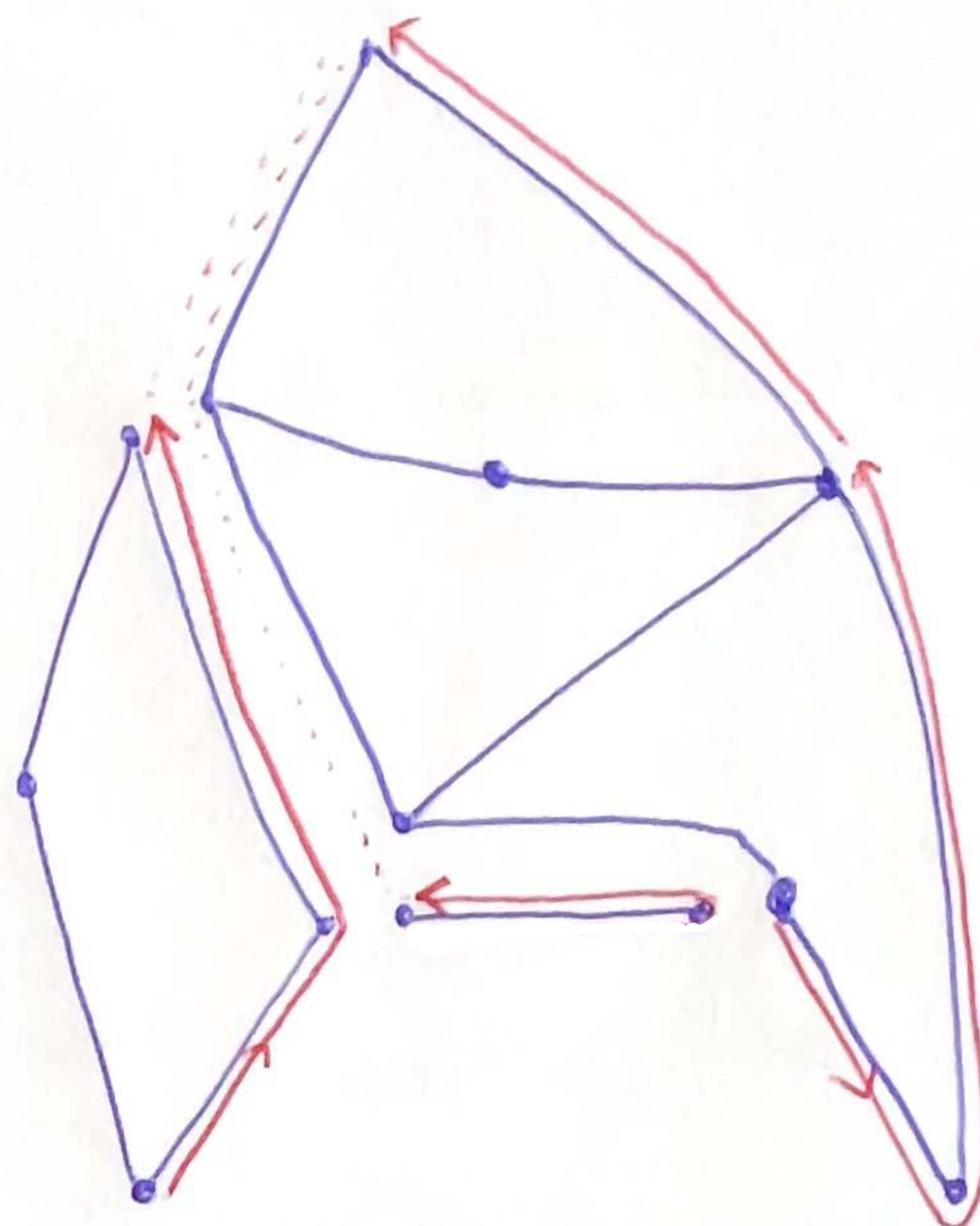
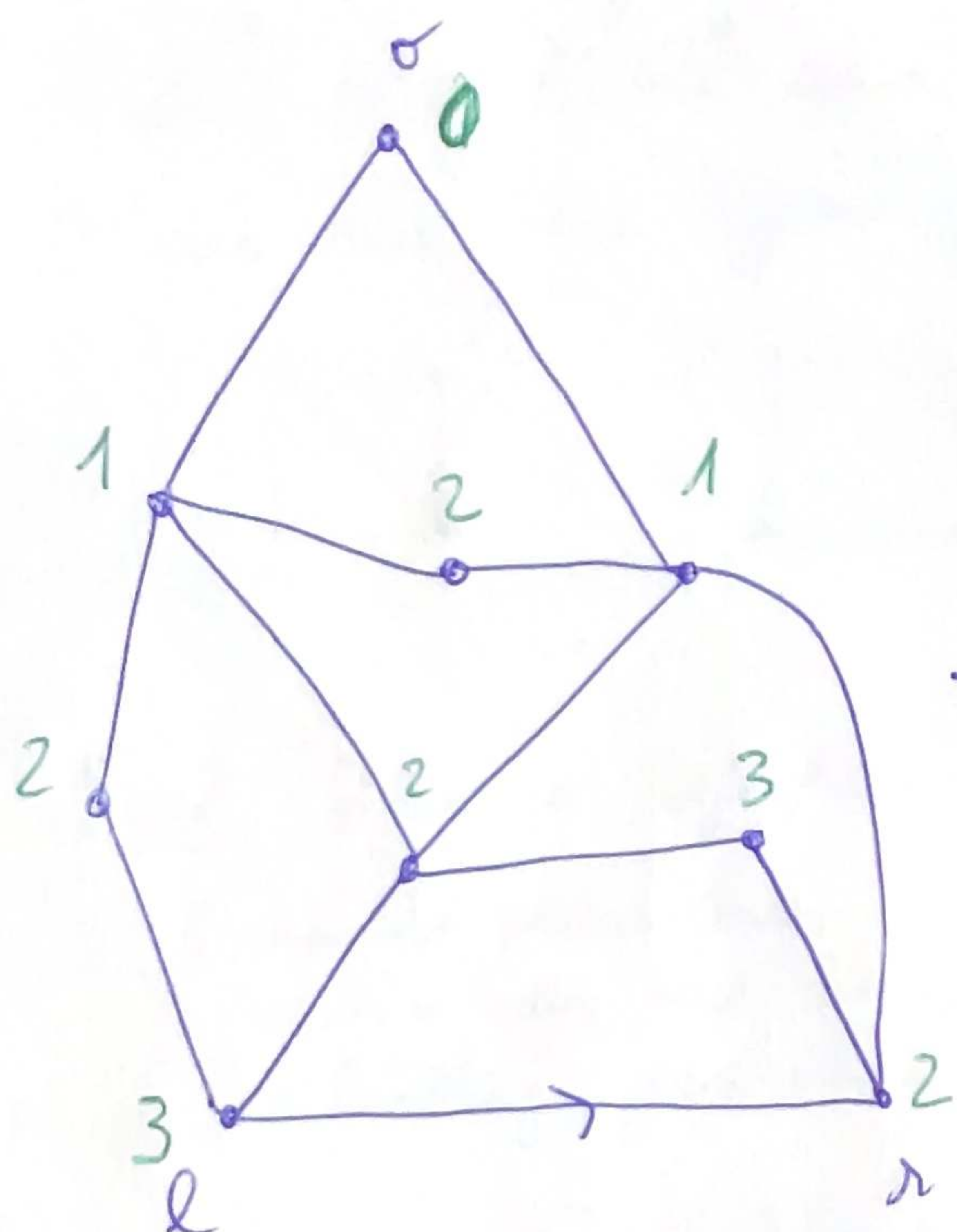
The g.f we wanted!

Finally: slices of type $R(z) \xleftrightarrow{bij}$ elements of $\mathcal{Q}_m$ where the root-edge is oriented $i+1 \rightarrow i$



$$\rightarrow |\mathcal{Q}_m| = (m+2)|\mathcal{Q}_m| = 2 \cdot 3^m \text{Cat}(m) !$$

→ generalizes almost directly to $2p$ -angulations
and in fact bipartite maps with all degrees controlled.
cf exercise session!!

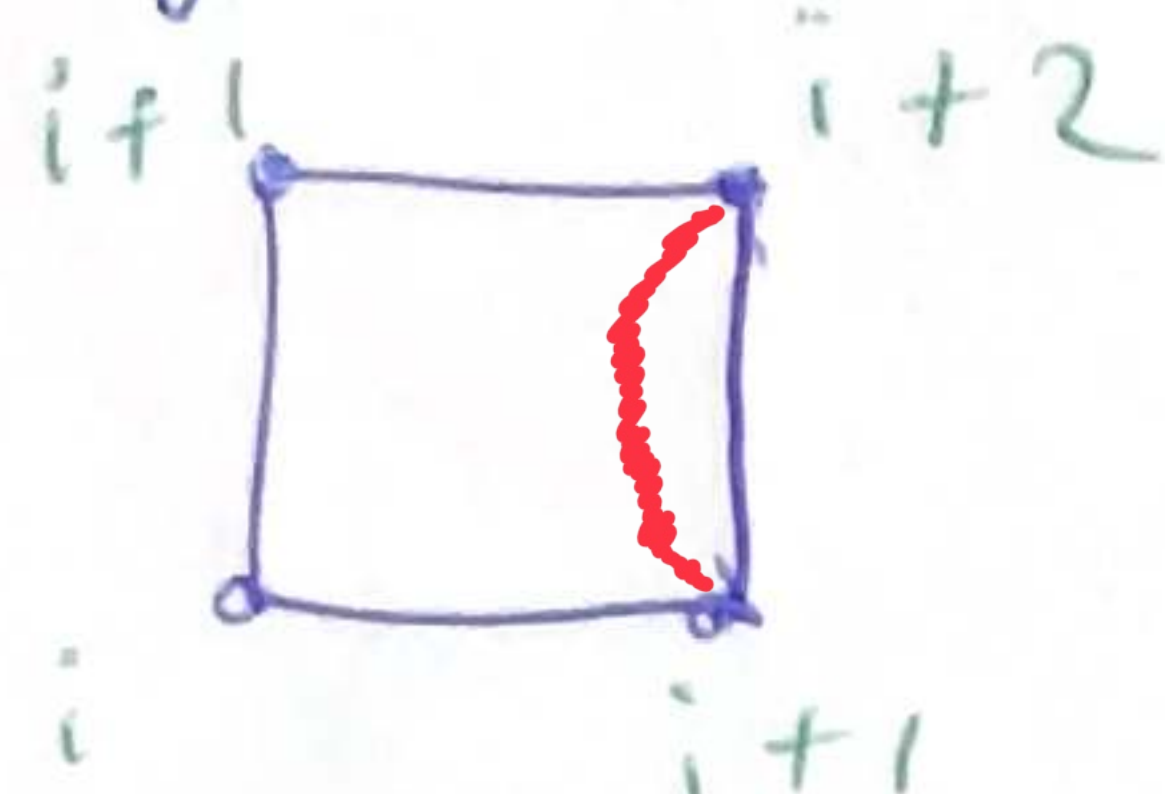
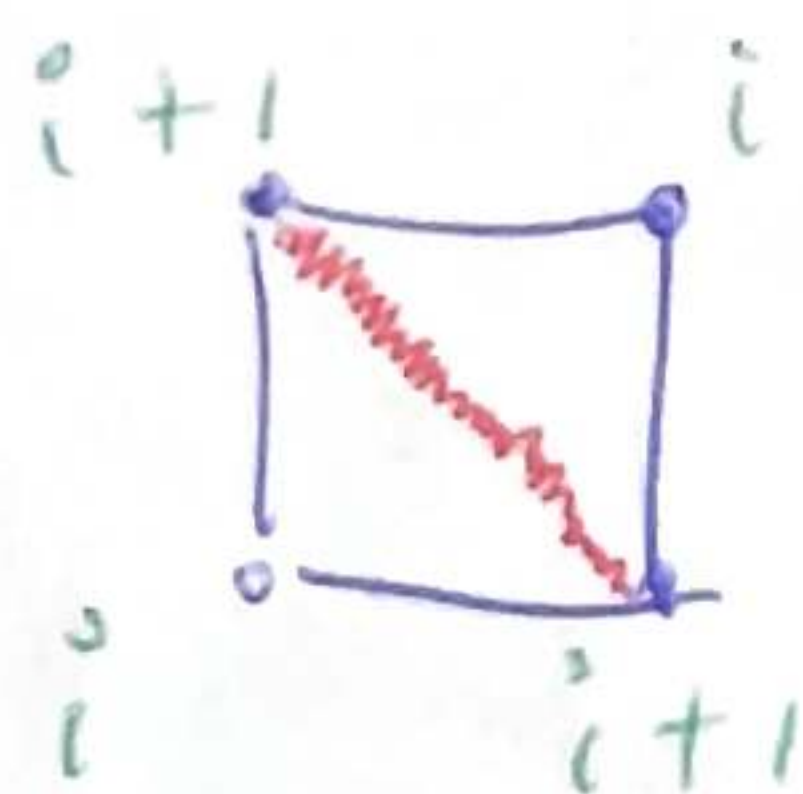
ex :

Tree interpretation (anterior-Schaeffer's bijection)

12

Take $(g, v) \in Q_n^*$ and label vertices by distance to v .

There are two types of faces:



Schaeffer rules: add red edge

Fact: This is a bijection:

{ rooted plane trees
with n edges, and 1-Lipschitz
vx labelling (global min = 1) }

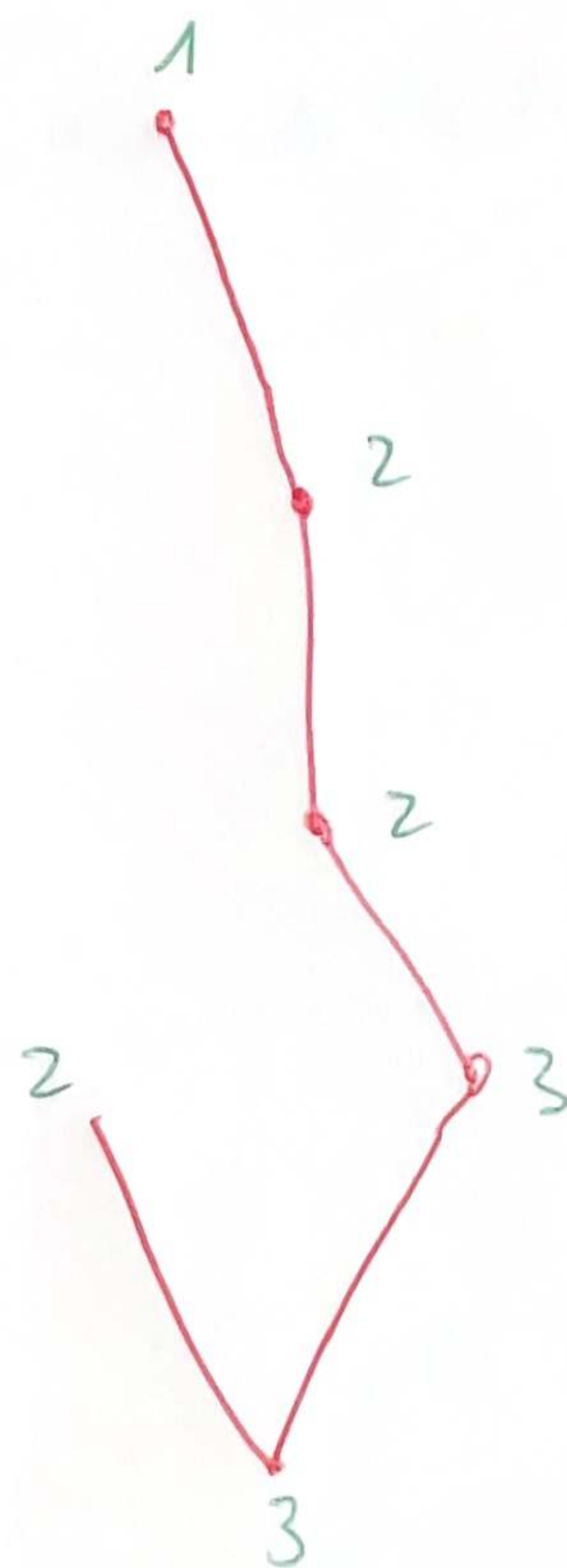
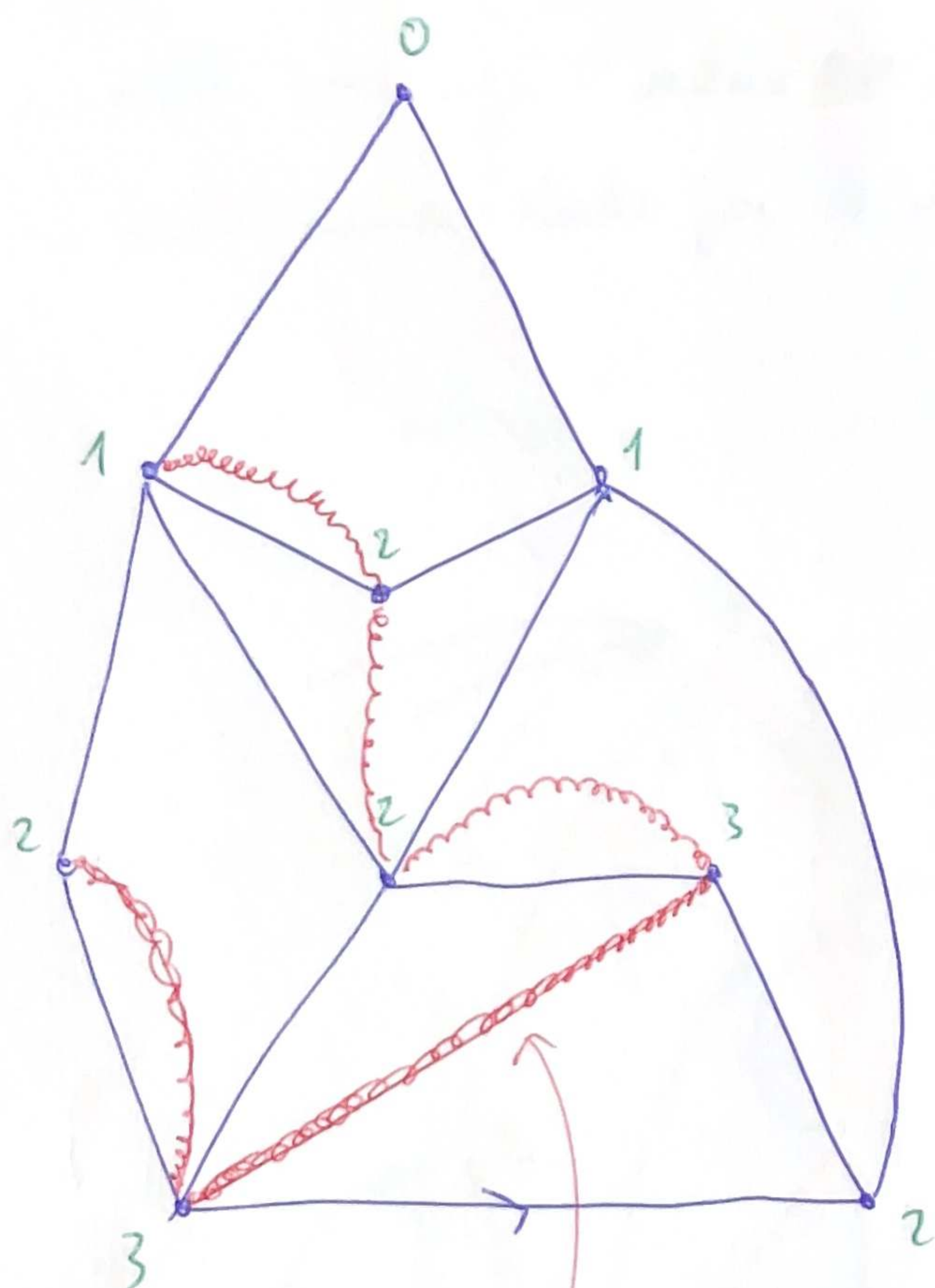
{ elements of Q_n^* with
root oriented $i+1 \rightarrow i$ }

$\tilde{3} \cdot \text{Cat}(n)$

$\longleftrightarrow \frac{1}{2} (n+2) Q_n$

In fact, this is the same as the slice construction:

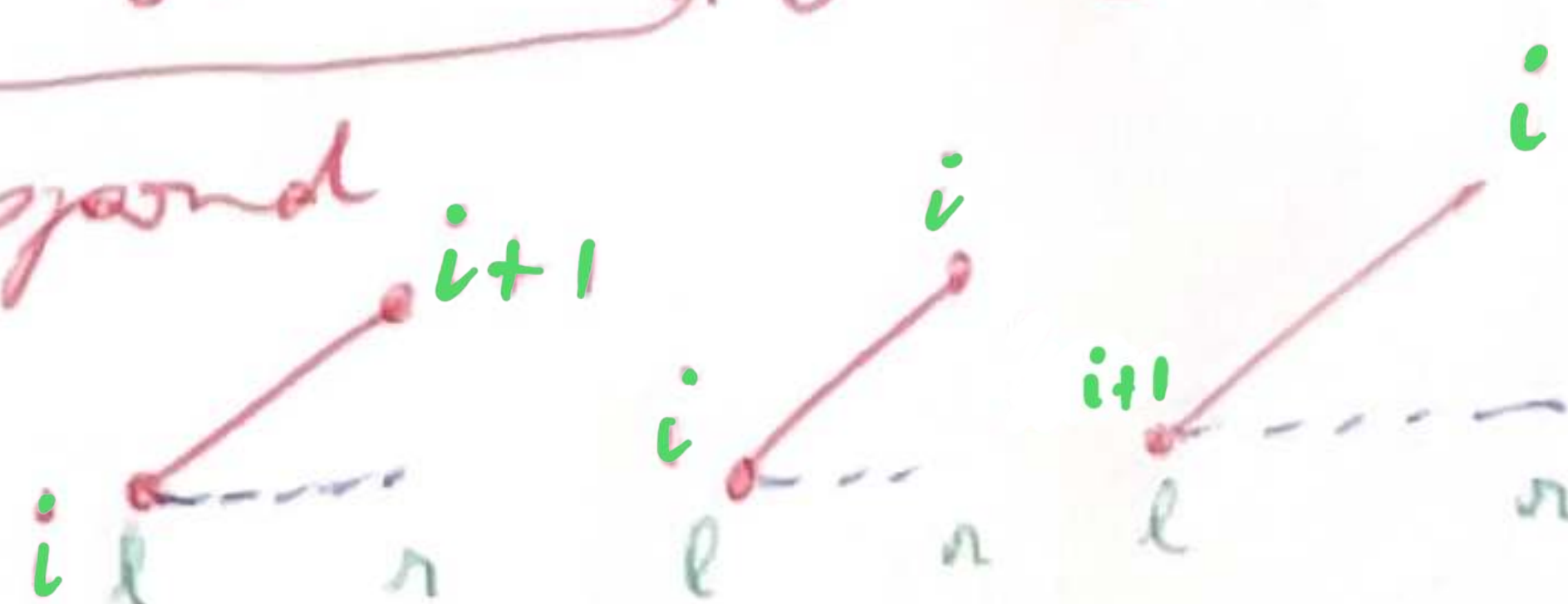
ex: (on a slice)



around the face, retains/links
only the two corners followed
by a descent (ie. the ones that
were giving rise ~~to~~ the
non-trivial " $R(z)$ slices".

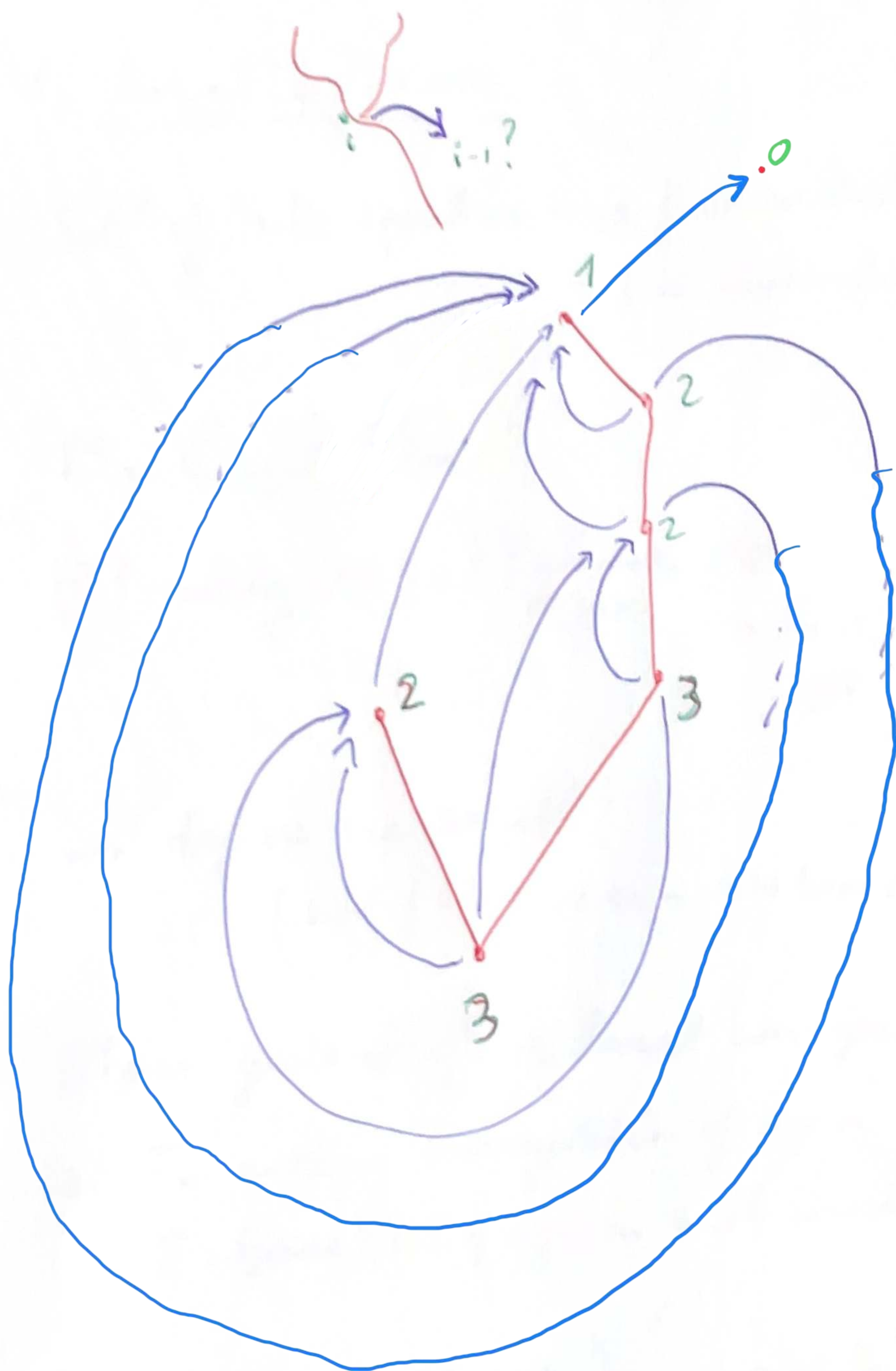
→ The recursive structure of this
tree = the recursive structure
of the slice equation $R(z) = 1 + \underline{3} z R(z)^2$

(and the three [♥] cases correspond
to the three types of edges



converse bijection:

- add missing "vertex $\overset{0}{\bullet}$ ".
- each corner i looks for its next $i-1$ around the tree:



IV Randomness:

Fundamental historical (but highly non-trivial!) consequences. Starting points of intense research areas.

1/ Local behaviour

Sol^o of Tutte equation $\rightarrow$ full control of root face degree (or root vertex, by duality).

M_n Euphr $\mathcal{M}_n$.

$P(\text{root dg}(M_n) = h) \xrightarrow{n \rightarrow \infty} p_h$ is some explicit finite sequence with exponential tail
 $p_h \sim c h^{+1/2} 2^{-h}$

$\rightarrow$ degrees are small!

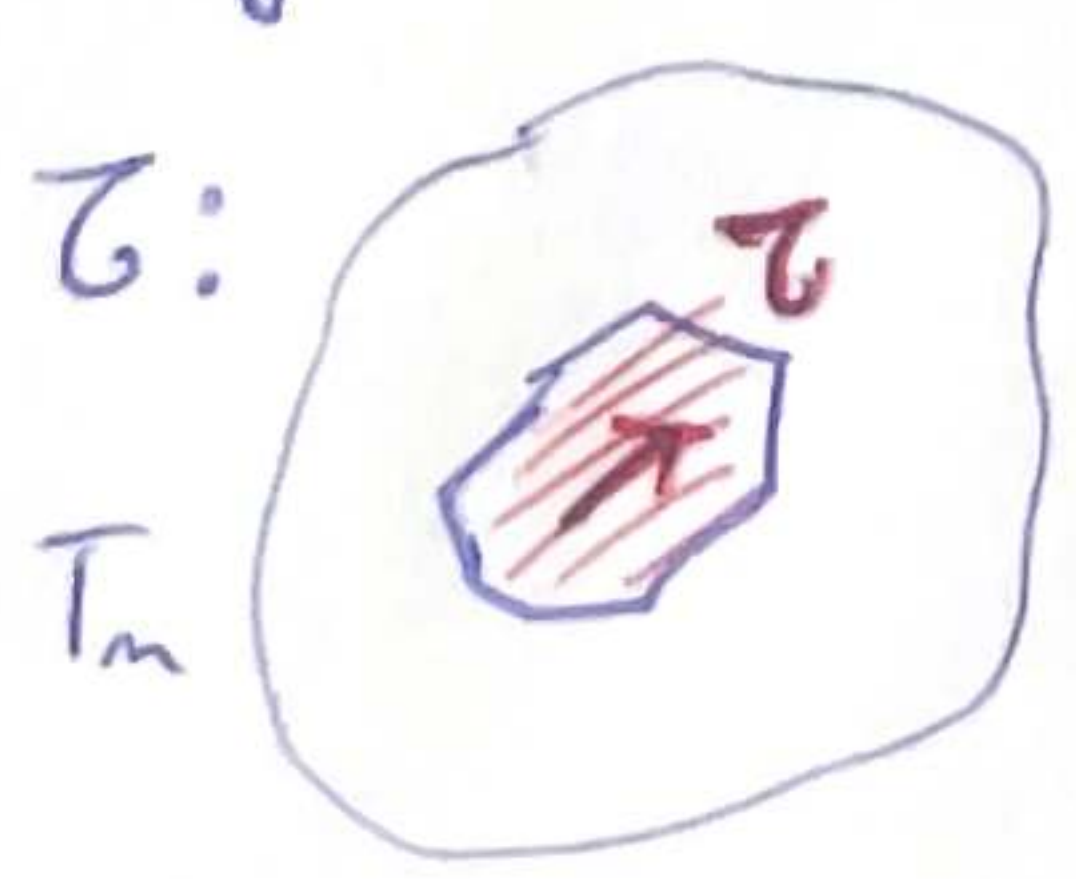
(kind of the same as in GW trees / plane trees).

More generally $\rightarrow$ local convergence.

Let T_n uniform triangulation of size n and

$\mathcal{T}$ = (fixed) tri^o of size m with boundary p .

$$P \left(\begin{array}{l} \text{neighbourhood of root in } T_n \\ \text{looks like } \mathcal{T}: \end{array} \right) = \frac{n \text{ of tri}^o (\text{size } n-m, \text{ perimeter } p)}{n \text{ of tri}^o (\text{size } n)}$$



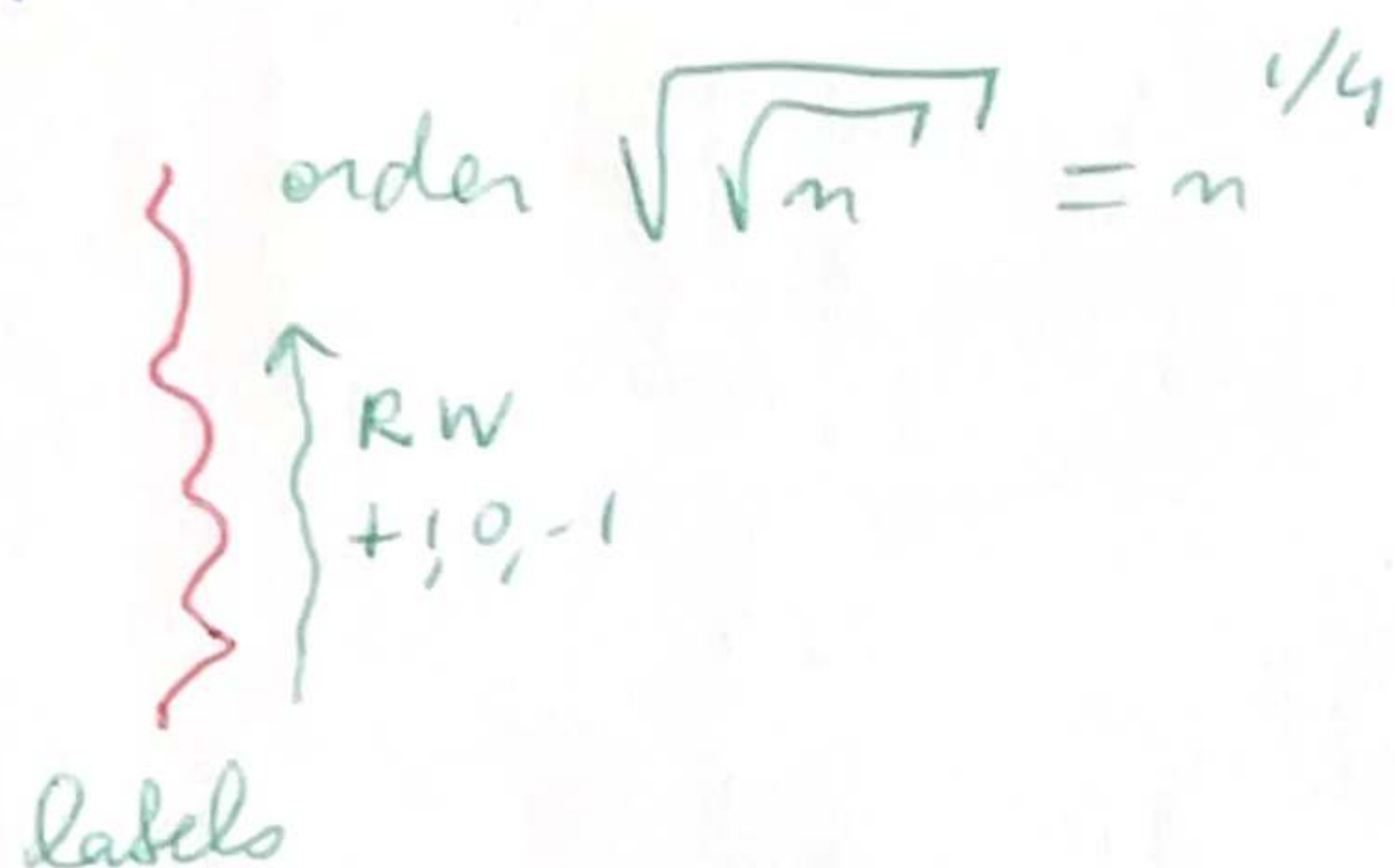
$\xrightarrow{n \rightarrow \infty}$ explicit nb ap.

[uses analog of Tutte's eq for triangulations. See exercise]

Notion of local limit, Uniform Infinite Random Triangulation (UIPT, Angel & Schram 2000's)

2/ Global behaviour

Schaeffer bijection $\Rightarrow$ distances to a random vertex in random quadrangulation Q_n are distributed like labels in a random 1-Lipschitz labelled random tree



Fact: $\text{diameter}(Q_n) = O_p(n^{1/4})$ when $n \rightarrow \infty$

More generally

$\frac{Q_n}{n^{1/4}} \xrightarrow[n \rightarrow \infty]{\text{some topology on metric spaces}}$

Brownian map

[Marchetti-Mukherjee, Le Gall, Micunovic, ...]

constructed by (difficult) continuous analogue of the Schaeffer bijection on continuous trees.

THE END

(Thanks!).